



Buffered Coscheduled MPI: A New Approach in the System Software Design for Large-Scale Parallel Computers

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Motivation



Hardware / OSs are glued together by System Software: Resource Management Communications Parallel Development and Debugging Tools Parallel File System Fault Tolerance





System software is a key factor to maximize usability, performance and scalability on large-scale systems!!!

Motivation

System software complexity due to multiple factors:

- Extremely complex global state
 - * Thousands of processes, threads, open files, pending messages, etc.
- Non-deterministic behavior
 - Inherent to computing systems ⇒ OS process scheduling
 - Induced by parallel applications ⇒ MPI_ANY_SOURCE
- Local OSs lack global awareness of parallel applications
 - ➤ Interferences with fine-grain synchronization operations ⇒ Non-scalable collective communication primitives.





Motivation

System software complexity due to multiple factors:

• Independent design of different components

➤ Redundancy of functionality ⇒ Communication protocols

- → Missing functionality ⇒ QoS user-level traffic / system-level traffic
- User-level applications rely on system software

System software performance/scalability impacts user-application performance/scalability







Target

- Simplifying design and implementation of the system software for large-scale computers
- Simplicity, performance, scalability

Approach

- Built atop a basic set of three primitives
- Global synchronization/scheduling

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Vision

- SIMD system running MIMD applications
 - (variable granularity in the order of hundreds of μ s)



Outline

- Motivation and Goals
- Introduction
- Core Primitives
- Design and Implementation
- Performance Evaluation
- Concluding remarks





Introduction

In this paper we make our point by implementing the communication layer: BCS MPI

- Built atop the three core primitives we proposed as the basics for all system software components (STORM is a resource manager implemented atop the same primitives)
- SIMD-like behavior: communications are synchronized and scheduled every few hundreds of microseconds
- Implemented almost entirely in the Network Interface Card
- Sacrifices neither performance nor scalability
- Paves the way to provide improved system utilization, traffic segregation, deterministic replay of user applications and system-level fault tolerance







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System software built atop three primitives

- Xfer-And-Signal
 - Transfer block of data to a set of nodes
 - Optionally signal local/remote event upon completion
- Compare-And-Write
 - Compare global variable on a set of nodes
 - Optionally write global variable on the same set of nodes
- Test-Event
 - Poll local event





System software built atop three primitives

• Xfer-And-Signal (QsNet):

- Node S transfers block of data to nodes D₁, D₂, D₃ and D₄
- Events triggered at source and destinations



System software built atop three primitives

• Compare-And-Write (QsNet):

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- Node S compares variable V on nodes D₁, D₂, D₃ and D₄



System software built atop three primitives

- Compare-And-Write (QsNet):
 - Node S compares variable V on nodes D₁, D₂, D₃ and D₄
 - Partial results are combined in the switches







Outline





Real-time communication scheduling



Global synchronization

- Strobe sent at regular intervals (time slices)
 - Compare-And-Write + Xfer-And-Signal (Master)
 - Test-Event (Slaves)
- All system activities are tightly coupled
- Global information is required to schedule resources, global synchronization facilitates the task but it is not enough

Global Scheduling

- Exchange of communication requirements
 - Xfer-And-Signal + Test-Event
- Communication scheduling
- Real transmission
 - Xfer-And-Signal + Test-Event





Implementation in the Network Interface Card

- Application processes interact with NIC threads
 - MPI primitive ⇒ Descriptor posted to the NIC
 - Communications are buffered
- Cooperative threads running in the NIC
 - Synchronize
 - Partial exchange of control information
 - Schedule communications
 - Perform real transmissions and reduce computations
- Computation/communication completely overlapped
 - Incoming messages do not generate interrupts
 - User processes do not need to poll for communication completion



Non-blocking primitives: MPI_lsend/lrecv







Blocking primitives: MPI_Send/Recv



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Global Synchronization/Scheduling Protocol

- Global Message Scheduling Phase
 - Microphases: Descriptor Exchange + Message Scheduling
- Message Transmission Phase:
 - Microphases: Point-to-point, Barrier and Broadcast, Reduce





Processes and Threads

- Synchronization: Strobe Sender + Strobe Receiver
- Scheduling: Buffer Sender, Buffer Receiver
- Transmission: DMA, Collective and Reduce Helpers







Simplicity

• Hierarchical design atop a single kernel module which implements the three core primitives

Determinism

- Synchronization facilitates resource scheduling
- Resource scheduling enforces reproducibility

Performance / Scalability

Hardware-supported primitives



BCS MPI implementation

QsNet hardware-supported multicast / events



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Concluding remarks

BCS MPI vs. Quadrics MPI

- Experimental Setup
 - crescendo cluster at LANL/CCS-3
 - 32 Dell 1550 compute nodes (two 1GHz P-III processors)
 - Quadrics QM-400 Elan3 NIC (1 RAIL)
 - RH 7.3 + Quadrics mods + qsnetlibs v1.5.0-0
 - Intel C/Fortran Compiler v5.0.1 (-O3)
 - accelerando cluster at LANL/CCS-3
 - 32 HP Server rx2600 (two 1GHz Itanium-II processors)
 - Two Quadrics QM-400 Elan3 NICs (2 RAILs)
 - RH 7.2 + Quadrics mods + qsnetlibs v1.5.0-0
 - Intel C/Fortran Compiler v7.1.17 (-O3)







BCS MPI vs. Quadrics MPI

- Experimental Setup
 - User-level implementation of BCS MPI
 - Scheduling parameters
 - > 500µs communication scheduling time slice (1 rail)
 - > 250µs communication scheduling time slice (2 rails)
 - Benchmarks and Applications
 - NPB (IS,EP,MG,CG,LU) Class C
 - SWEEP3D 50x50x50
 - SAGE timing.input





Benchmarks and Applications (IA32)

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SAGE - timing.input (IA32)

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Blocking vs Non-blocking SWEEP3D (IA32) MPI_Send/Recv \Rightarrow MPI_Isend/Irecv + MPI_Waitall



Blocking vs Non-blocking SWEEP3D (IA64) MPI_Send/Recv \Rightarrow MPI_Isend/Irecv + MPI_Waitall



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BCS MPI: Similar Performance with a Simplified Design





Concluding Remarks

New approach to the design of system software

- Kernel module which implements the basics of most system software components (STORM) ⇒ Simplicity
- Hardware-supported primitives \Rightarrow Performance / Scalability

Prototype implementation of BCS MPI

- Global synchronization of all system activities
- Global scheduling of communications
- Implemented almost entirely in the NIC
- Global state optimization vs latency/bandwidth optimization

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Promising preliminary results with real apps

BCS MPI can compete with a production-level MPI



Future Work

Kernel-level implementation of BCS-MPI

• User-level solution is already working with negligible performance degradation

Improved system utilization

• Scheduling multiple jobs

QoS for different types of traffic

• Scheduling messages may provide traffic segregation

Deterministic replay of MPI programs

• Ordered resource scheduling may enforce reproducibility



Transparent fault tolerance

• BCS MPI simplifies the state of the machine







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